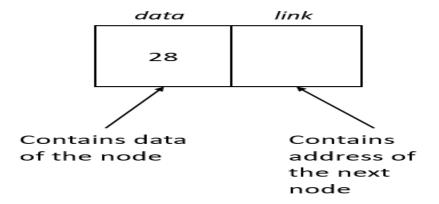
Linear / Linked List

Types of linear/linked lists:

- 1) Single Linked List
- 2) Doubly Linked List
- 3) Circular Linked List

Single Linked List:

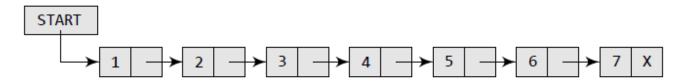
- A linked list is a linear collection of data elements. These data elements are called nodes.
- In single linked list every node contains two fields, data field and link field -a pointer to the next node/address of next node.



The node in a single linked list is declared as

```
struct node
{
int data;
struct node *next;
};
```

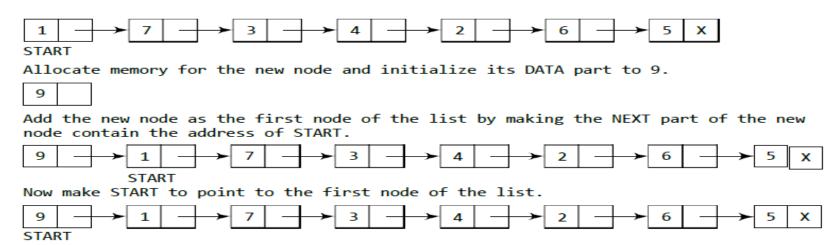
➤ The last node address field in the single linked list contains NULL.



- > Operations performed on a single linked list:
 - I. Insertion
 - II. Deletion
 - III. Searching
 - IV. Traversing

Insertion:

Case 1: The new node is inserted at the beginning



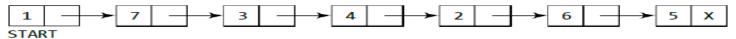
Step 4: SET NEW_NODE -> DATA = VAL

Step 5: SET NEW_NODE -> NEXT = START

Step 6: SET START = NEW_NODE

Step 7: EXIT

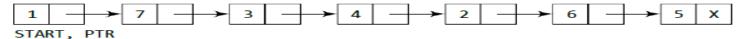
Case 2: The new node is inserted at the end



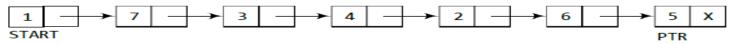
Allocate memory for the new node and initialize its DATA part to 9 and NEXT part to NULL.

9 X

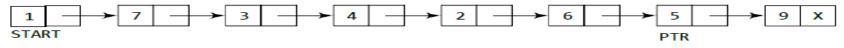
Take a pointer variable PTR which points to START.



Move PTR so that it points to the last node of the list.



Add the new node after the node pointed by PTR. This is done by storing the address of the new node in the NEXT part of PTR.



- Step 4: SET NEW_NODE > DATA = VAL
- Step 5: SET NEW_NODE > NEXT = NULL
- Step 6: SET PTR = START
- Step 7: Repeat Step 8 while PTR > NEXT != NULL
- Step 8: SET PTR = PTR > NEXT

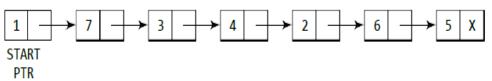
 [END OF LOOP]
- Step 9: SET PTR > NEXT = NEW_NODE
- Step 10: EXIT

Case 3: The new node is inserted after a given node

Allocate memory for the new node and initialize its DATA part to 9.

9

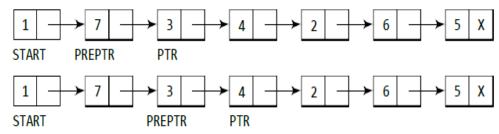
Take two pointer variables PTR and PREPTR and initialize them with START so that START, PTR, and PREPTR point to the first node of the list.



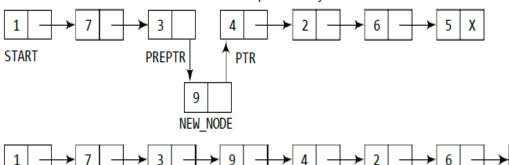
PREPTR

START

Move PTR and PREPTR until the DATA part of PREPTR = value of the node after which insertion has to be done. PREPTR will always point to the node just before PTR.



Add the new node in between the nodes pointed by PREPTR and PTR.



Step 4: SET NEW_NODE - > DATA = VAL

Step 5: SET PTR = START

Step 6: SET PREPTR = PTR

Step 7: Repeat Steps 8 and 9 while PREPTR - > DATA != NUM

Step 8: SET PREPTR = PTR

Step 9: SET PTR = PTR - > NEXT

[END OF LOOP]

Step 10: PREPTR - > NEXT = NEW_NODE

Step 11: SET NEW_NODE - > NEXT = PTR

Step 12: EXIT

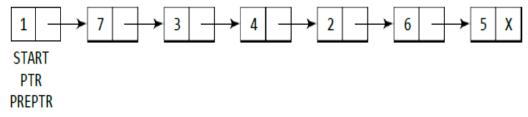
Case 4: The new node is inserted before a given node

Allocate memory for the new node and initialize its DATA part to 9.

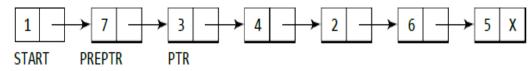
9

START

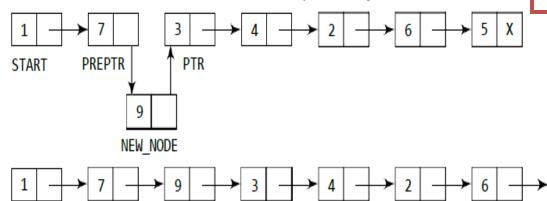
Initialize PREPTR and PTR to the START node.



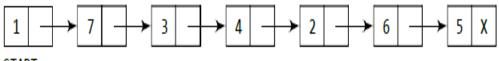
Move PTR and PREPTR until the DATA part of PTR = value of the node before which insertion has to be done. PREPTR will always point to the node just before PTR.



Insert the new node in between the nodes pointed by PREPTR and PTR.

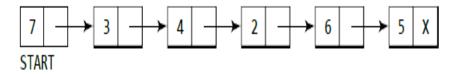


Deletion : Case 1: The first node is deleted



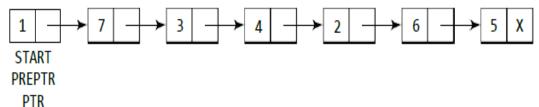
START

Make START to point to the next node in sequence.

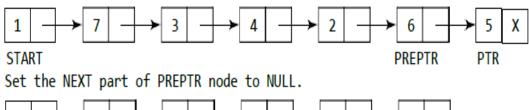


Case 2: The last node is deleted

Take pointer variables PTR and PREPTR which initially point to START.



Move PTR and PREPTR such that NEXT part of PTR = NULL. PREPTR always points to the node just before the node pointed by PTR.



```
Step 1: IF START = NULL

Write UNDERFLOW

Go to Step 5

[END OF IF]

Step 2: SET PTR = START

Step 3: SET START = START -> NEXT

Step 4: FREE PTR

Step 5: EXIT
```

```
Step 1: IF START = NULL

Write UNDERFLOW

Go to Step 8

[END OF IF]

Step 2: SET PTR = START

Step 3: Repeat Steps 4 and 5 while PTR -> NEXT != NULL

Step 4: SET PREPTR = PTR

Step 5:SET PTR = PTR -> NEXT [END

OF LOOP]

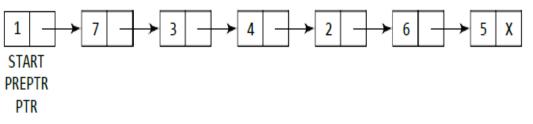
Step 6: SET PREPTR -> NEXT = NULL

Step 7: FREE PTR

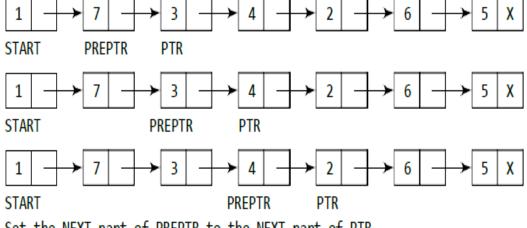
Step 8: EXIT
```

Case 3: The node after a given node is deleted

Take pointer variables PTR and PREPTR which initially point to START.



Move PREPTR and PTR such that PREPTR points to the node containing VAL and PTR points to the succeeding node.



Set the NEXT part of PREPTR to the NEXT part of PTR.

START

```
Step 1: IF START = NULL

Write UNDERFLOW

Go to Step 10

[END OF IF]

Step 2: SET PTR = START

Step 3: SET PREPTR = PTR

Step 4: Repeat Steps 5 and 6 while PREPTR -> DATA != NUM

Step 5: SET PREPTR = PTR

Step 6:SET PTR = PTR -> NEXT [END

OF LOOP]

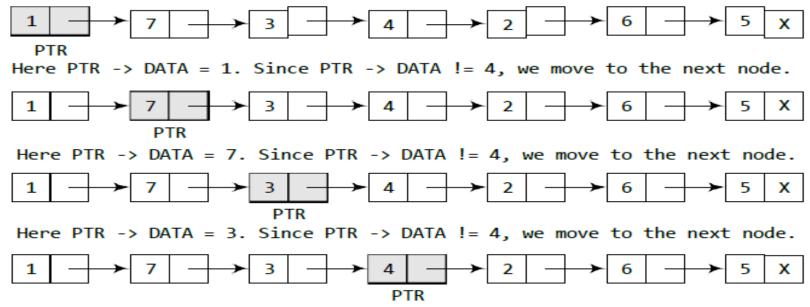
Step 7: SET TEMP = PTR

Step 8: SET PREPTR -> NEXT = PTR -> NEXT

Step 9: FREE TEMP

Step 10: EXIT
```

Searching:



Here PTR -> DATA = 4. Since PTR -> DATA = 4, POS = PTR. POS now stores the address of the node that contains VAL

```
Step 1: [INITIALIZE] SET PTR = START
Step 2: Repeat Step 3 while PTR != NULL
Step 3: IF VAL = PTR -> DATA

SET POS = PTR
Go To Step 5

ELSE

SET PTR = PTR -> NEXT

[END OF IF]

[END OF LOOP]
Step 4: SET POS = NULL
Step 5: FXIT
```

> Traversing:

Traversing a linked list means accessing the nodes of the list in order to perform some processing on them.

```
Step 1: [INITIALIZE] SET PTR = START

Step 2: Repeat Steps 3 and 4 while PTR != NULL

Step 3: Apply Process to PTR -> DATA

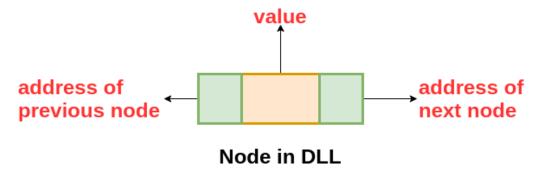
Step 4:SET PTR = PTR -> NEXT [END OF

LOOP]

Step 5: EXIT
```

Doubly Linked List

- A doubly linked list or a two-way linked list is a more complex type of linked list which contains a pointer to the next as well as the previous node in the sequence.
- It is a collection of node, in which each node will contain three fields- a pointer to the previous node, data, a pointer to the next node.



The declaration of node in double linked list is given as

```
struct node
{
    struct node *prev;
    int data;
    struct node *next;
};
```

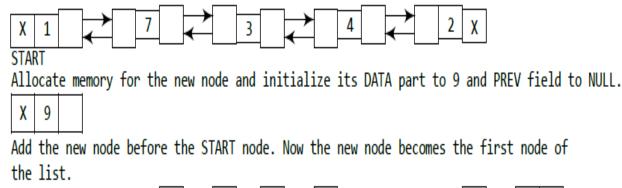
- The PREV field of the first node and the NEXT field of the last node will contain NULL.
- The NEXT field is used to traverse the list in forward direction and PREV field is used to traverse the list in backward direction.

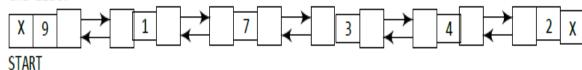
Operation on Doubly linked list are

- I. Insertion
- II. Deletion
- III. Searching
- IV. Traversing

> Insertion:

Case 1: The new node is inserted at the beginning.





```
SET NEW_NODE -> DATA = VAL

SET NEW_NODE -> PREV = NULL

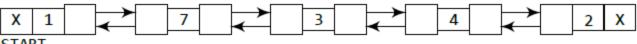
SET NEW_NODE -> NEXT = START

SET START -> PREV = NEW_NODE

SET START = NEW_NODE

EXIT
```

Case 2: The new node is inserted at the end.

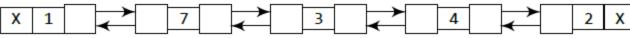


START

Allocate memory for the new node and initialize its DATA part to 9 and its NEXT field to NULL.

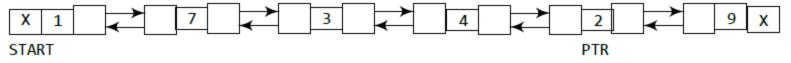


Take a pointer variable PTR and make it point to the first node of the list.



START, PTR

Move PTR so that it points to the last node of the list. Add the new node after the node pointed by PTR.



```
Step 4: SET NEW_NODE -> DATA = VAL
```

Step 5: SET NEW_NODE -> NEXT = NULL

Step 6: SET PTR = START

Step 7: Repeat Step 8 while PTR -> NEXT != NULL

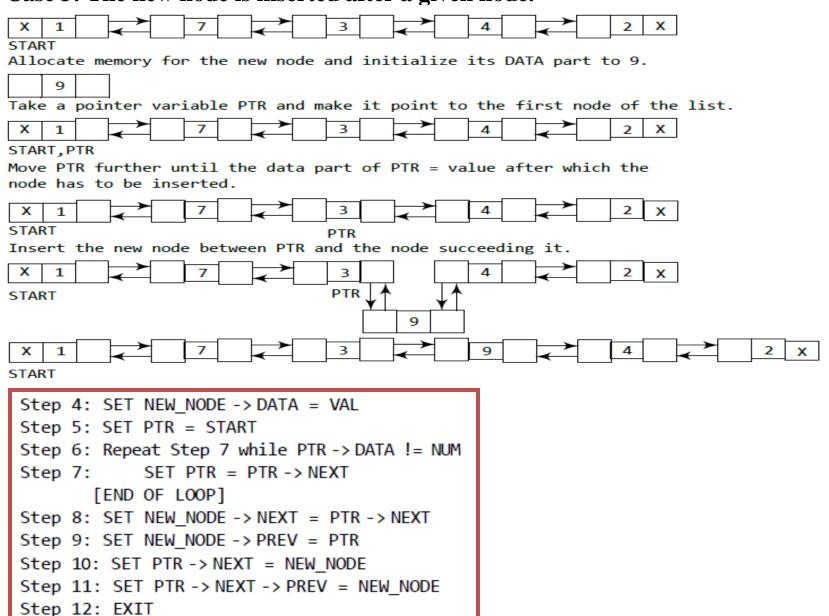
Step 8: SET PTR = PTR -> NEXT [END OF LOOP]

Step 9: SET PTR -> NEXT = NEW_NODE

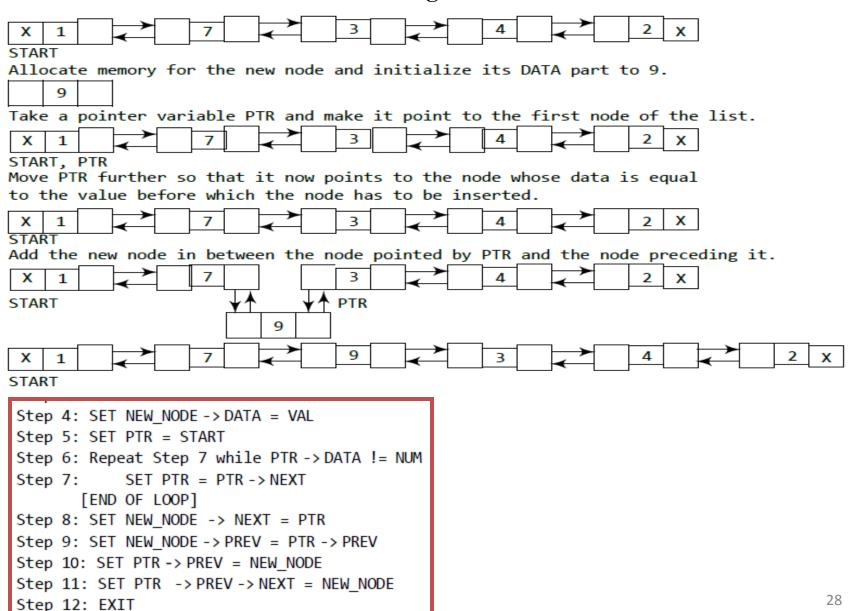
Step 10: SET NEW_NODE -> PREV = PTR

Step 11: EXIT

Case 3: The new node is inserted after a given node.

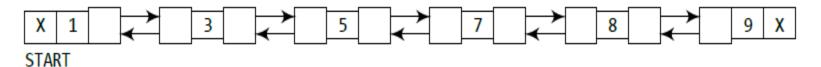


Case 4: The new node is inserted before a given node.

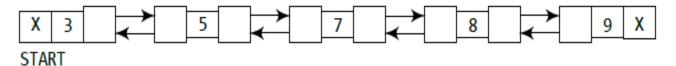


Deletion:

Case 1: The first node is deleted.



Free the memory occupied by the first node of the list and make the second node of the list as the START node.



```
Step 1: IF START = NULL

Write UNDERFLOW

Go to Step 6

[END OF IF]

Step 2: SET PTR = START

Step 3: SET START = START -> NEXT

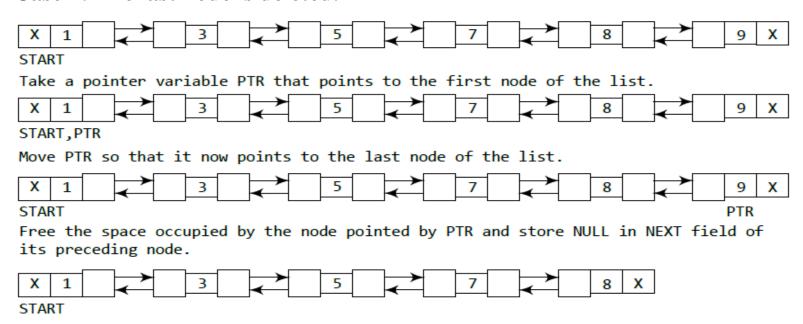
Step 4: SET START -> PREV = NULL

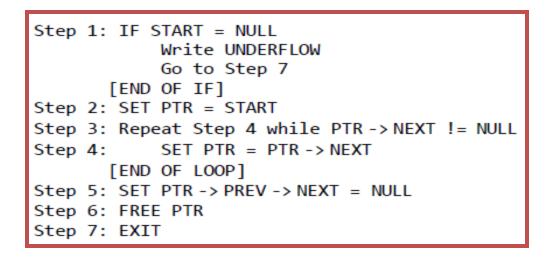
Step 5: FREE PTR

Step 6: EXIT
```

Deletion:

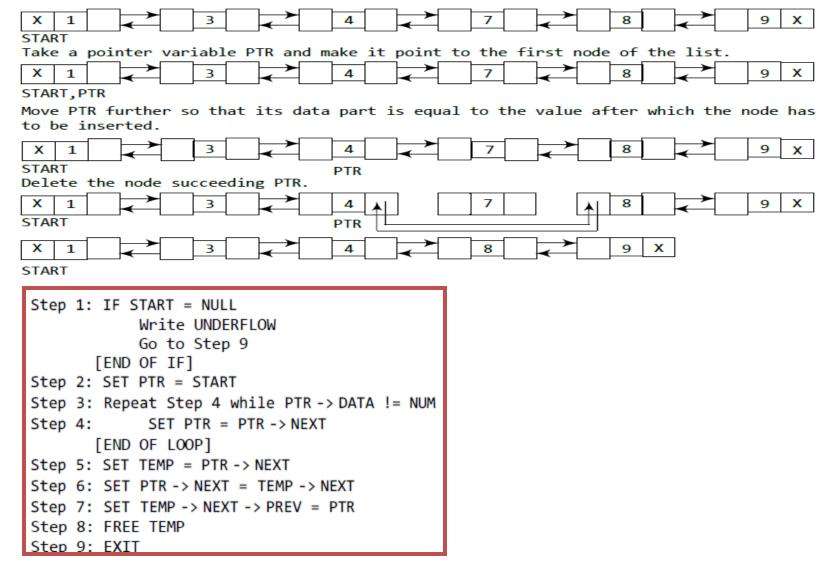
Case 2: The last node is deleted.





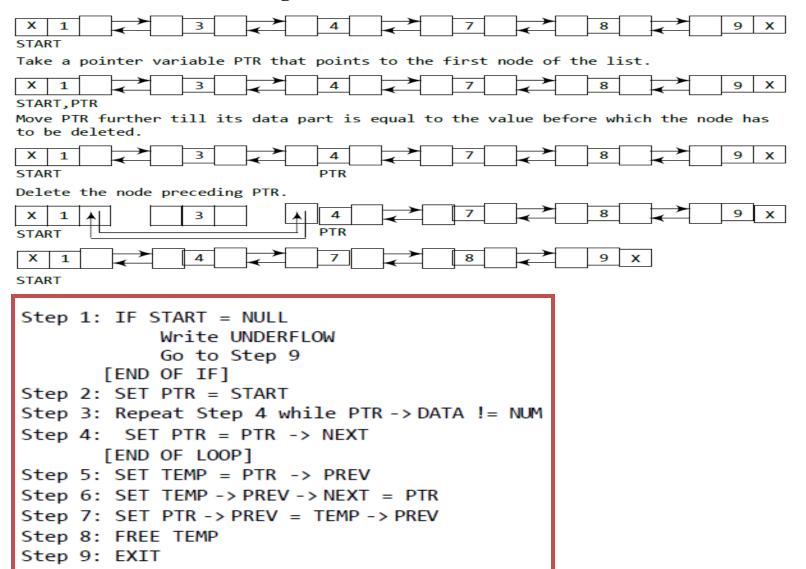
Deletion:

Case 3: The node after a given node is deleted.



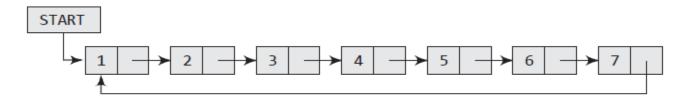
Deletion:

Case 4: The node before a given node is deleted.



Circular Linked List

- Circular Linked List Types:
 - I. Circular Single Linked List
 - II. Circular Doubly Linked List
- Circular Single Linked List:
 - In a circular single linked list, the last node contains a pointer to the first node of the list i,e the last node address field contains the address of the first node.

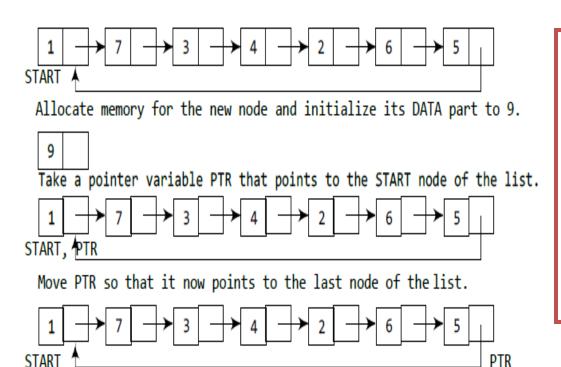


- Operation on circular single linked list are:
 - I. Insertion
 - II. Deletion
 - III. Searching
 - IV. Traversing

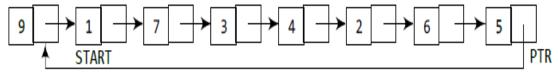
Circular Single Linked List.....

Insertion:

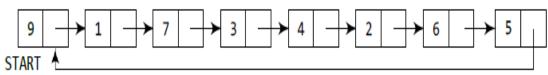
Case 1: The new node is inserted at the beginning of the circular linked list.



Add the new node in between PTR and START.



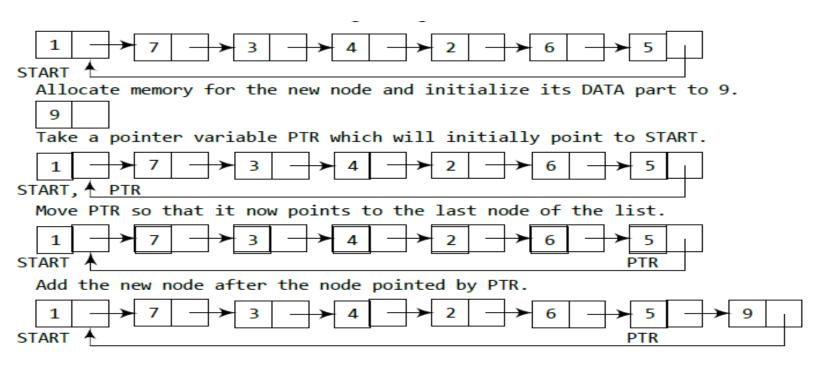
Make START point to the new node.



Circular Single Linked List.....

Insertion:

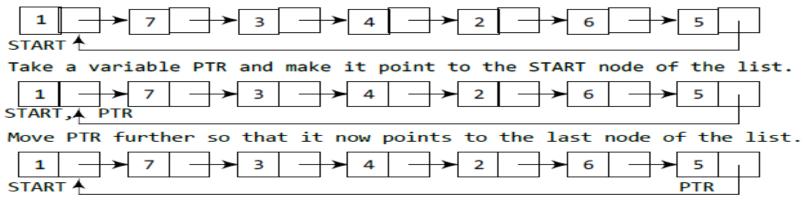
Case 2: The new node is inserted at the end of the circular linked list.



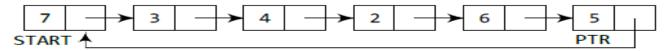
Circular Single Linked List....

Deletion:

Case 1: The first node is deleted.



The NEXT part of PTR is made to point to the second node of the list and the memory of the first node is freed. The second node becomes the first node of the list.



```
Step 1: IF START = NULL

Write UNDERFLOW

Go to Step 8

[END OF IF]

Step 2: SET PTR = START

Step 3: Repeat Step 4 while PTR -> NEXT != START

Step 4: SET PTR = PTR -> NEXT

[END OF LOOP]

Step 5: SET PTR -> NEXT = START -> NEXT

Step 6: FREE START

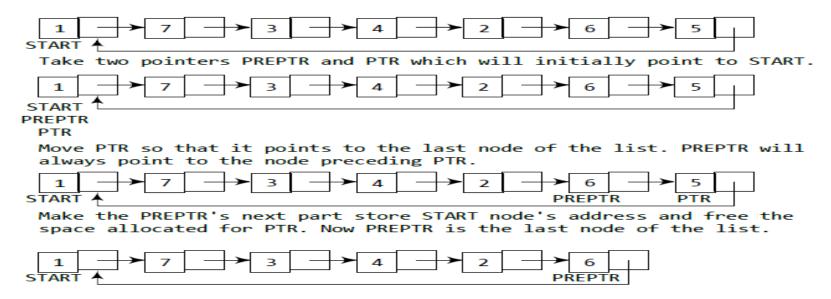
Step 7: SET START = PTR -> NEXT

Step 8: EXIT
```

Circular Single Linked List....

Deletion:

Case 2: The last node is deleted.



```
Step 1: IF START = NULL

Write UNDERFLOW

Go to Step 8

[END OF IF]

Step 2: SET PTR = START

Step 3: Repeat Steps 4 and 5 while PTR -> NEXT != START

Step 4: SET PREPTR = PTR

Step 5: SET PTR = PTR -> NEXT

[END OF LOOP]

Step 6: SET PREPTR -> NEXT = START

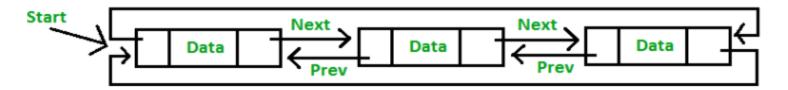
Step 7: FREE PTR

Step 8: EXIT
```

Circular Doubly Linked List

Circular Doubly Linked List:

- Circular doubly linked list doesn't contain NULL in any of the node.
- The last node NEXT field of the list contains the address of the first node of the list.
- The first node PREV field of the list contain address of the last node of the list.

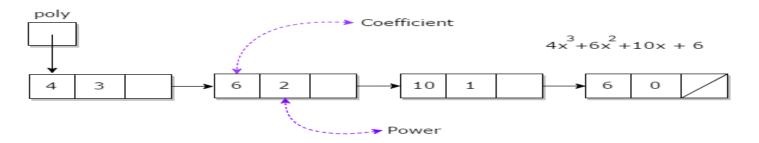


Operations:

- I. Insertion
- II. Deletion
- III. Searching
- IV. Traversing
- Implementation is more complex than other linked lists.

Applications of Linked Lists

Polynomial Representation



- > Implementation of different data structures like stack, queues etc.
- Dynamic Memory allocation.
- Used in Image viewer-previous and next images are linked, hence we can access by using next and previous buttons.
- In Operating System, all the running applications are kept in a circular linked list and the OS gives a fixed time slot to all for running.
- Used in web browser to access the next and previous web pages while browsing.
